An MPI Halo-Cell Implementation for Zero-Copy Abstraction

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Jean-Baptiste Besnard (1), Allen Malony (2), Sameer Shende (1), Marc Pérache (3), Patrick Carribault (3) and Julien Jaeger (3)

ParaTools SAS, Bruyères-le-Châtel
 ParaTools Inc, Eugene USA
 CEA, DAM, DIF F91297 Arpajon France

jbbesnard@paratools.fr



Introduction (1/3)

HPC machines are rapidly shifting to higher concurrency

- Now gathering millions of cores
- Intra-node parallelism is rapidly increasing (several hundred threads) (Xeon Phi / KNL)
- This with a smaller memory per thread

It is well aknowledged that applications will have to evolve in order to take advantage of such architectures MPI + X being often refered to as a potential solution.



Introduction (2/3)

But what does it mean...

► What is this X ?

Distributed Memory	Shared-Memory	Accelerators	Logical Address Spaces
MPI and optimized intra-node communications	OpenMP Cilk, TBB Pthreads, 	GPUs, FPGAs	PGAS, DSM

There are several alternatives: MPI + OpenMP, MPI + GPU, MPI + PGAS,



Introduction (3/3)

But what does it mean...

► Why MPI is not sufficient ? Why do we need this X ?

In our paper, we propose to model this limitation when considering domain-splitting in distributed memory context

We show that distributed memory poses problems of:

- Memory due to domain replication
- Communication overhead and therefore scalability

Then, we propose an MPI level abstraction solving these issues for domain splitting by providing the advantages of shared-memory programming.

Domain Splitting (1/2)



We consider the case where computation is done over a distributed domain (often as a stencil) creating dependencies between cells structured as a mesh

—> This covers a wide range of applications (not all)

Domain Splitting (2/2)

It is common knowledge for all MPI programmers that such domain spitting requires halo/ghost cells on local domain boundaries.

Is it possible to provide a simple model of the halo-cells ? What is the performance impact for common topologies ?

Yes (first part of our paper)



Domain Splitting on four processes Paratools

Halo-Cell Model (1/6)

To derive this model, we considered wrapped-around meshes (tori) instead of regular ones in order to have a regular mesh layout (no border effect).



Halo-Cell Model (2/6)



These regular topologies are nonetheless completely representative of unwrapped ones dealing with the level of connectivity between distributed areas. 8

Halo-Cell Model (3/6)

n: Number of cells
C: Number of halo layers
d: Mesh dimension
I: Characteristic length of the topology

$$l(n,d) = n^{\frac{1}{d}}$$



Paratools

$$N_g(n,d) = (l(n,d) + 2C)^d - l(n,d)^d = (n^{\frac{1}{d}} + 2C)^d - n$$

 $\!$ Subtract a mesh without halo-cells to a mesh with a characteristic length increased of 2C. $\!$

Halo-Cell Model (4/6)



2 * 3 (processes) * 1 (layer) = 6 halo cells



Halo-Cell Model (5/6)



$$4pC(\sqrt{\frac{n}{p}}+C)$$

2D

= 4.4.1(sqrt(16/4) + 1)= 4.4.1(sqrt(4) + 1) = 4.4.3

= 48 = (4 * 12) halo cells

Halo-Cell Model (6/6)



2D (two layers)

$$4pC(\sqrt{\frac{n}{p}}+C)$$

= 4.4.2(sqrt(16 / 4) + 2)

= 4.4.2(sqrt(4) + 2)

= 4.4.2.4

= 128 = (4 * 32) halo cells



Halo-Cells and Performance (1/4) $S(n,p) = \frac{s(n)}{\frac{s(n)}{p}} + comm(n,p)$

Starting from the well-known speedup equation, it can be seen that strong-scaling speedup is bounded by communications which are directly linked to the number of halo-cells.

—> Computation time should be much larger than communication time. There should be more local cells than halo cells with a complex computation.
 However, this ratio changes with p (strong scaling)

Halo-Cells and Performance (2/4)

If we now consider the weak-scaling model, we have n/p which is a constant as is the ghost cell ratio.

Communication cost has then to be independent of the number of processes, in order to allow weakscaling. Which is true for regular decomposition?



Halo-Cells and Performance (3/4)

	1D	2D	3D
$\boxed{n(r,p,C)}$	$\frac{2pC}{r}$	$\frac{4pC^2}{r^2}(r+2+2\sqrt{1+r})$	$\frac{2pC^3}{r^3}(r^2+3r(1+r)^{\frac{2}{3}}+6r(1+r)^{\frac{1}{3}})$
			$+9[r+1+(1+r)^{\frac{2}{3}}+(1+r)^{\frac{1}{3}}])$
$\boxed{n(1\%, p, 1)}$	200p	$1.61 \times 10^{5} p$	$2.18 \times 10^{8} p$
n(1%,p,2)	400p	$6.43 \times 10^5 p$	$1.74 \times 10^{9} p$
$\fbox{n(1\%,p,3)}$	600p	$1.44 \times 10^{6} p$	$5.89 \times 10^9 p$
n(10%, p, 1)	20p	1679p	$2.37 \times 10^5 p$
$\fbox{n(10\%,p,2)}$	40p	6716p	$1.90 imes 10^6 p$
$\fbox{n(10\%,p,3)}$	60 <i>p</i>	15111 <i>p</i>	$6.42 \times 10^6 p$
n(50%, p, 1)	4p	79p	2639p
$\fbox{n(50\%,p,2)}$	8p	317 <i>p</i>	21177 <i>p</i>
$\boxed{n(50\%, p, 3)}$	16p	712p	71272 <i>p</i>

When doing weak-scaling, it is desirable to limit the ghostcell ratio in order to completely hide communication costs.

However, memory per thread is decreasing:

In 3D, if you want 1% of ghost cells with one layer, you need 1.64 GB of memory (for 8 bytes cells). Compare it to the 34 MB / Thread on a Xeon Phi.



Halo-Cells and Performance (4/4)



Hybrid Approach

Intra-node parallelism is then a direct way of reducing the ghost cell ratio and then improving scalability by overcoming the per thread memory limitation.

Reducing communication cost

 Limiting ghost-cell memory overhead while freeing memory for computation (hiding comms)



MPI Optimized Intra-Node Messaging

A lot of work has been done to optimize intra-node communications:

- SHM memory segments
- KNEM kernel module
- Or since Linux 3.2 Cross Memory Attach (CMA)
- Direct copy in thread-based MPI
- It is even possible to use the HCA to emit RDMA

Such approaches efficiently reduce node-local communication cost but do not reduce/remove the memory associated with halo cells which still has to be duplicated.



MPI Halo

We propose a Halo Cell abstraction providing the advantages of shared-memory models while remaining close to MPI semantics:

- Transparent use of larger memory areas
- Removal of memory duplications between tasks on the same node
- Removal of node-local communications (no copies Zero copy)
- Support for computation outside of node boundaries (no mixing)

MPI Halo Principle (1/4)



Classical Ghost Cell Approach With Copies

When doing a stencil, most applications use two meshes, one for « t » and another for « t+1 », approach required due to the spatial-dependency between cells.

MPI Halo Principle (2/4)



What if local cells (located on the same node) could be resolved as local pointers — no copies would be required.

The source mesh being accessed in read-only is not necessary to duplicate data.

MPI Halo Principle (3/4)



Pointer exchanges allow mesh-switching

MPI Halo Principle (4/4)



Illustration of both inter-node and intra-node exchanges with MPI-Halo cells.



MPI Halo Example (1D splitting) (1/2)

```
/*---- Initialization (Done once) */
MPI Halo local left, local right, left, right;
/* Name Cells and provide Layout */
MPIX_Halo_cell_init( &local_left, "Local Left" , MPI_INT, 1024 );
MPIX_Halo_cell_init( &local_right, "Local Right", MPI_INT, 1024);
MPIX Halo cell init( &left, "Remote Right", MPI INT, 1024);
MPIX_Halo_cell_init( &right, "Remote Left", MPI_INT, 1024);
/* Bind Cells */
MPI_Halo_ex ex;
MPIX_Halo_exchange_init( &ex );
MPIX_Halo_cell_bind_local( ex, local_left );
MPIX_Halo_cell_bind_local( ex, local_right );
MPIX_Halo_cell_bind_remote( ex, right, right_process, "Local Left" );
MPIX_Halo_cell_bind_remote( ex, left, left_process, "Local Right" );
/* Generate Communications */
MPIX_Halo_exchange_commit( ex );
```



MPI Halo Example (1D splitting) (2/2)

/*---- Compute Loop (Called at each time-step)*/
while(compute)

{

}

```
/* Register local cell data */
MPIX Halo cell set( local left, mesh );
MPIX Halo cell set( local right, right coll( mesh ) );
/* Start asynchronous communications */
MPIX Halo iexchange( ex );
/* ... Compute mesh center ... */
MPIX_Halo_iexchange_wait( ex );
/* Retrieve Ghost arrays */
int * left ghost, * right ghost;
MPIX_Halo_cell_get( left, (void **)&left_ghost );
MPIX Halo cell get( right, (void **)&right ghost );
/* ... Compute mesh boundaries ... */
/* Swap Meshes */
Mesh * tmp = mesh;
mesh = oldmesh;
oldmesh = tmp;
```

MPI Halo Interface

MPI_Halo:

- Automatic buffer abstraction (local or remote)
- Can be set to a value when local
- A pointer can be retrieved when remote
- Supports MPI data-types (packing abstraction)

MPI_Halo_ex:

- Build the communication scheme between MPI_Halo
- Buffers are named (no abstract offsets)
- An error is reported if the remote is not present
- No offset is passed to communication calls

Boundaries have to be handled as particular case

Copy can still be forced when the remote is modified



MPI Halo Performance Results (1/3)

Our test-case was the convolution of a 5616x3744 RBG image implemented in OpenMP, MPI-Halo, also forcing buffer allocation to behave like the classical ghost-cell approach. We tested this benchmark with various convolution kernel sizes.

Our MPI-Halo interface has been implemented in the MPC runtime which is a thread-based MPI, making node-level exchanges trivial (shared-memory). Nothing prevents the MPI Halo model to be ported to process-based MPI supposing a previous memory registration.



MPI Halo Performance Results (2/3)

- OpenMP (7x7) OpenMP (3x3) — MPI Halo and MPI Copy (11x11) $\vdash \times \dashv$ OpenMP (11x11) ---
- MPI Halo and MPI Copy (3x3)
- MPI Halo and MPI Copy (7x7) \mapsto



Task Count

Computation Time

MPI Halo Performance Results (3/3)



(b) Memory Usage (Resident)

Memory Usage

arato

Conclusion

Halo-Cell Model:

- Introduced a model of the halo-cell ratio
- Explained that scaling was highly impacted by this ratio
- Shown that distributed memory was hitting the perthread memory barrier, encouraging hybrid models to achieve better ghost-cell ratios particularly for higher dimensions (3D with several layers).

MPI_Halo:

- Proposed an MPI based solution to the domain decomposition issue we exposed (buffer aliasing)
- Allows a clear definition of a communication scheme with static validation of buffer matching (size, name)
- Consistent with inter-node parallelism (unlike OpenMP)
 Parallelism (unlike OpenMP)

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